

## MATH 245 S17, Exam 2 Solutions

1. Carefully define the following terms: free variable, predicate, counterexample, Left-to-Right Principle.

A free variable is a variable that has not been bound, drawn from some domain. A predicate is a collection of propositions, indexed by one or more free variables, each drawn from its domain. A counterexample is a specific domain element chosen to make a predicate false. The Left-to-Right Principle states that variables are bound from left to right.

2. Carefully define the following terms: Uniqueness Proof Theorem, Proof by Contradiction Theorem, Proof by Induction Theorem, well-ordered set.

The Uniqueness Proof theorem states: there is at most one domain element satisfying predicate  $P$  if  $\forall x, y \in D, P(x) \wedge P(y) \rightarrow x = y$ . The Proof by Contradiction theorem states: For propositions  $p, q$ , if  $(p \wedge \neg q) \equiv F$ , then  $p \rightarrow q$  is true. The Proof by Induction theorem states: To prove  $\forall x \in \mathbb{N}, P(x)$ , we prove (a)  $P(1)$  is true; and (b)  $\forall x \in \mathbb{N}, P(x) \rightarrow P(x+1)$ .

3. Simplify  $\neg(\exists x, \forall y, \forall z, (x < y) \rightarrow (x < z))$  as much as possible (i.e. where nothing is negated). Do not prove or disprove this statement.

$\forall x, \exists y, \exists z, (x < y) \wedge (x \geq z)$ , or (nicer)  $\forall x, \exists y, \exists z, z \leq x < y$ .

4. Recall that  $\mathbb{R} \setminus \mathbb{Q}$  is the set of irrational numbers. Let  $a \in \mathbb{R} \setminus \mathbb{Q}, b \in \mathbb{Q}$ . Use proof by contradiction to prove that  $a + b \in \mathbb{R} \setminus \mathbb{Q}$ .

Because  $b \in \mathbb{Q}$ , there are integers  $m, n$  with  $n \neq 0$  and  $b = \frac{m}{n}$ . Now, assume by way of contradiction that  $a + b \in \mathbb{Q}$ . Then there are integers  $s, t$  with  $t \neq 0$  and  $a + b = \frac{s}{t}$ . We calculate  $a = (a + b) - b = \frac{s}{t} - \frac{m}{n} = \frac{sn - mt}{nt}$ . Now,  $sn - mt, nt$  are integers, and  $nt \neq 0$ , so  $a \in \mathbb{Q}$ . This is a contradiction.

5. Prove or disprove:  $\forall x \in \mathbb{R}, \forall y \in \mathbb{R}, (x < y) \rightarrow \lceil x \rceil \leq \lfloor y \rfloor$ .

The statement is false. To prove this, we need to prove  $\exists x \in \mathbb{R}, \exists y \in \mathbb{R}, (x < y) \wedge (\lceil x \rceil > \lfloor y \rfloor)$ . Take  $x = 0.3, y = 0.4$ . We have  $x < y$  but  $\lceil x \rceil = 1 > 0 = \lfloor y \rfloor$ . Many other  $x, y$  are possible.

6. Let  $n \in \mathbb{Z}$ . Use the Division Algorithm to prove that  $\frac{(n-1)(n+2)}{2} \in \mathbb{Z}$ .

Apply DA to get integers  $q, r$  with  $n = 2q + r$ , and  $0 \leq r < 2$ . We now have two cases. If  $r = 0$ , then  $\frac{(n-1)(n+2)}{2} = \frac{(n-1)(2q+0+2)}{2} = (n-1)q \in \mathbb{Z}$ . If instead  $r = 1$ , then  $\frac{(n-1)(n+2)}{2} = \frac{(2q+1-1)(n+2)}{2} = q(n+2) \in \mathbb{Z}$ .

7. Recall the Fibonacci numbers given by  $F_0 = 0, F_1 = 1, F_n = F_{n-1} + F_{n-2}$  (for  $n \geq 2$ ). Prove that for all  $n \in \mathbb{N}_0, F_{n+2} = 1 + \sum_{i=0}^n F_i$ .

Base case  $n = 0$ :  $F_2 = 1 = 1 + F_0 + F_1$ . Inductive case: Let  $n \in \mathbb{N}_0$  and assume that  $F_{n+2} = 1 + \sum_{i=0}^n F_i$ . Add  $F_{n+1}$  to both sides:  $F_{n+3} = F_{n+1} + F_{n+2} = F_{n+1} + 1 + \sum_{i=0}^n F_i = 1 + \sum_{i=0}^{n+1} F_i$ . Hence  $F_{n+3} = 1 + \sum_{i=0}^{n+1} F_i$ .

8. Let  $x \in \mathbb{R}$ . Prove that  $\lfloor x \rfloor$  exists. That is, prove  $\exists n \in \mathbb{Z}, n \leq x < n + 1$ .

Let  $S$  be the set of all integers less than or equal to  $x$ . This is a nonempty set, with an upper bound ( $x$ ), so by the Maximum Element Induction theorem, there is some maximum element  $n \in S$ . Since  $n \in S, n \leq x$ . We now prove  $x < n + 1$ . Assume, by way of contradiction, that  $x \geq n + 1$ . But then  $n + 1 \in S$ , and  $n + 1 > n$ , a contradiction since  $n$  was a maximum. Hence  $n \leq x < n + 1$ .

9. Use induction to prove  $\forall n \in \mathbb{N}, \frac{(2n)!}{n!n!} \geq 2^n$ .

Base case  $n = 1$ :  $\frac{2!}{1!1!} = 2 \geq 2^1$ . Inductive case: Let  $n \in \mathbb{N}$ , and assume that  $\frac{(2n)!}{n!n!} \geq 2^n$ . Multiply both sides by  $\frac{(2n+1)(2n+2)}{(n+1)(n+1)}$  to get  $\frac{(2(n+1))!}{(n+1)!(n+1)!} = \frac{(2n)!}{n!n!} \frac{(2n+1)(2n+2)}{(n+1)(n+1)} \geq 2^n \frac{(2n+1)(2n+2)}{(n+1)(n+1)} = 2^n \frac{(2n+1)2(n+1)}{(n+1)(n+1)} = 2^{n+1} \frac{2n+1}{n+1} \geq 2^{n+1}$ . Hence  $\frac{(2(n+1))!}{(n+1)!(n+1)!} \geq 2^{n+1}$ .

10. Let  $\mathbb{R}^+$  denote the positive real numbers. Prove that  $\forall a \in \mathbb{R}^+, \exists b \in \mathbb{R}^+, \forall x \in \mathbb{R}, |x - 2| < b \rightarrow |3x - 6| < a$ .

Let  $a \in \mathbb{R}^+$  be arbitrary. Choose  $b = \frac{a}{3}$ . Now, let  $x \in \mathbb{R}$  with  $|x - 2| < b$ . We have  $|x - 2| < b = \frac{a}{3}$ . Multiplying both sides by 3, we get  $|3x - 6| = 3|x - 2| < a$ . Hence  $|3x - 6| < a$ . This proves that  $\lim_{x \rightarrow 2} 3x = 6$ ; to learn much more like this, take Math 534A.